

Pattern-avoiding ascent sequences

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Definition

An **ascent sequence** is a string $x_1 \cdots x_n$ of non-negative integers such that:

- ▶ $x_1 = 0$
- ▶ $x_n \leq 1 + \text{asc}(x_1 \cdots x_{n-1})$ for $n \geq 2$

\mathcal{A}_n is the set of ascent sequences of length n

$$\mathcal{A}_2 = \{00, 01\}$$

More examples: 01234, 01013

$$\mathcal{A}_3 = \{000, 001, 010, 011, 012\}$$

Non-example: 01024

Ascent Sequences

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Non-example: 01024

Theorem

$|\mathcal{A}_n|$ is the n th Fishburn number (OEIS A022493).

$$\sum_{n \geq 0} |\mathcal{A}_n| x^n = \sum_{n \geq 0} \prod_{i=1}^n (1 - (1-x)^i)$$

Definition

The **reduction** of $x = x_1 \cdots x_n$, $\text{red}(x)$, is the string obtained by replacing the i th smallest digits of x with $i - 1$.

Example: $\text{red}(273772) = 021220$

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Pattern containment/avoidance

$a = a_1 \cdots a_n$ **contains** $\sigma = \sigma_1 \cdots \sigma_m$ iff there exist $1 \leq i_1 < i_2 < \cdots < i_m \leq n$ such that $\text{red}(a_{i_1} a_{i_2} \cdots a_{i_m}) = \sigma$.

$$a_B(n) = |\{a \in \mathcal{A}_n \mid a \text{ avoids } B\}|$$

001010345 contains 012, 000, 1102; avoids 210.

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Goal

Determine $a_B(n)$ for many of choices of B .

Previous Work

- ▶ Duncan & Steingrímsson (2011)

Pattern b	$a_b(n)$	OEIS
001, 010 011, 012	2^{n-1}	A000079
102 0102, 0112	$(3^{n-1} + 1)/2$	A007051
101, 021 0101	$\frac{1}{n+1} \binom{2n}{n}$	A000108

- ▶ Mansour and Shattuck (2014)
Callan, Mansour and Shattuck (2014)

Patterns B	$a_B(n)$	OEIS
1012	$\sum_{\ell=0}^{n-1} \binom{n-1}{\ell} C_\ell$	A007317
0123	ogf: $\frac{1-4x+3x^2}{1-5x+6x^2-x^3}$	A080937
8 pairs of length 4 patterns	$\frac{1}{n+1} \binom{2n}{n}$	A000108

Other sequences (Baxter & P.)

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Patterns B	OEIS	$a_B(n)$
000,011	A000027	n
011,100	A000124	$\binom{n}{2} + 1$
001,210	A000125	$\binom{n}{3} + n$
000,101	A001006	M_n
000,001	A000045	F_{n+1}
001,100	A000071	$F_{n+2} - 1$
101,110	A001519	F_{2n-1}
100,101	A025242	(Generalized Catalan)
021,102	A116702	$ \mathcal{S}_n(123, 3241) $
102,120	A005183	$ \mathcal{S}_n(132, 4312) $
101,120	A116703	$ \mathcal{S}_n(231, 4123) $
201,210	A007317	$\sum_{\ell=0}^{n-1} \binom{n-1}{\ell} C_\ell$

Theorem (P.)

$$a_{201,210}(n) = \sum_{\ell=0}^{n-1} \binom{n-1}{\ell} C_\ell.$$

Theorem (Mansour & Shattuck)

$$a_{1012}(n) = \sum_{\ell=0}^{n-1} \binom{n-1}{\ell} C_\ell.$$

Conjecture (Duncan & Steingrímsson)

$$a_{0021}(n) = \sum_{\ell=0}^{n-1} \binom{n-1}{\ell} C_\ell.$$

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Conjecture (Duncan & Steingrímsson) / Theorem (P.)

$$a_{0021}(n) = \sum_{\ell=0}^{n-1} \binom{n-1}{\ell} C_\ell.$$

Definition

Patterns σ and ρ are **Wilf-equivalent** if $a_\sigma(n) = a_\rho(n)$ for $n \geq 1$. In this case, write: $\sigma \sim \rho$.

Example: $00 \sim 01$.

$a_{00}(n) = 1$ (the strictly increasing sequence)

$a_{01}(n) = 1$ (the all zeros sequence).

Definition

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Example: $00 \sim 01$.

$a_{00}(n) = 1$ (the strictly increasing sequence)

$a_{01}(n) = 1$ (the all zeros sequence).

Known from Duncan/Steingrímsson: All possible Wilf equivalences of length at most 4 are:

$$00 \sim 01$$

$$10 \sim 001 \sim 010 \sim 011 \sim 012$$

$$102 \sim 0102 \sim 0112$$

$$101 \sim 021 \sim 0101 \sim 0012$$

$$0021 \sim 1012$$

Theorem (P.)

$$a_{201,210}(n) = \sum_{\ell=0}^{n-1} \binom{n-1}{\ell} C_\ell.$$

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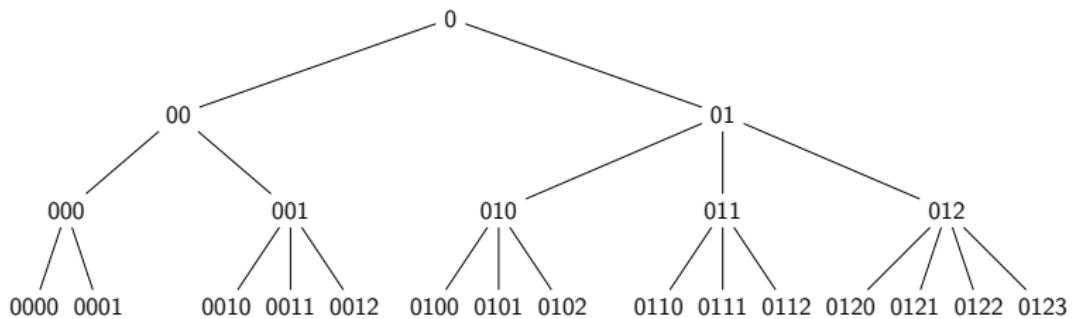
Conjecture (Duncan & Steingrímsson) / Theorem (P.)

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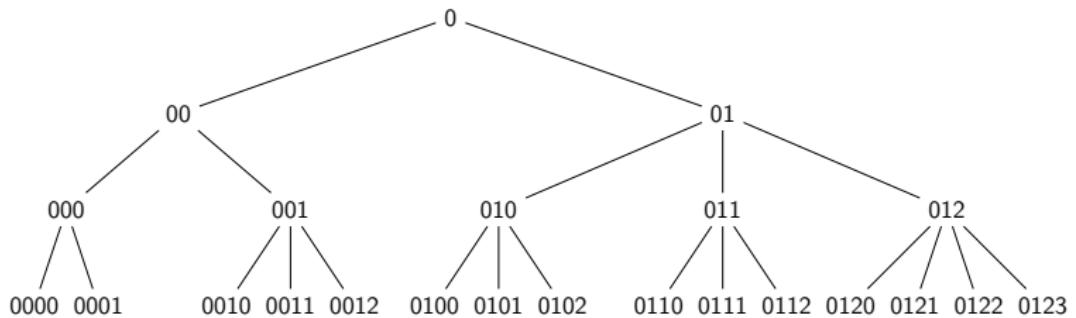
Proof scribble:

generating tree → recurrence → system of functional
equations → experimental solution → plug in for catalytic
variables

Generating Tree

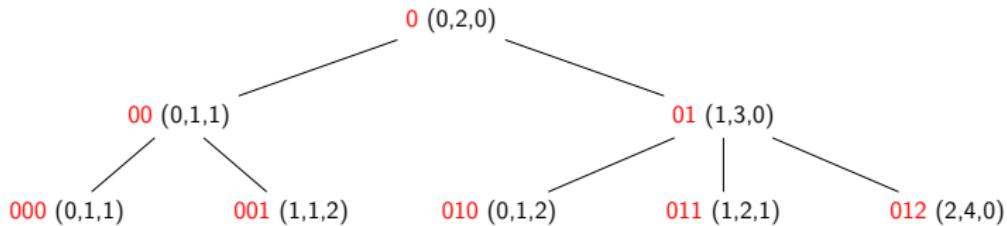


Generating Tree

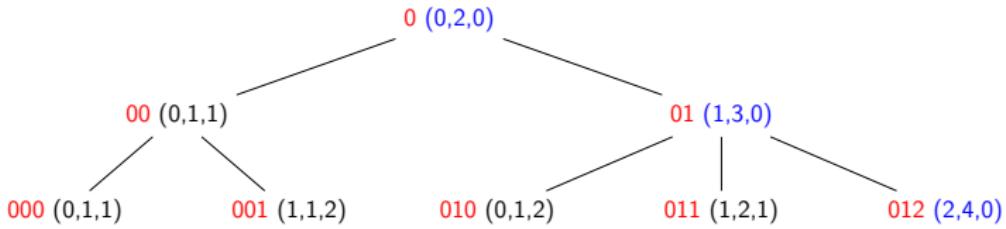


Idea: Replace a with an ordered triple of statistics on a .

Generating Tree

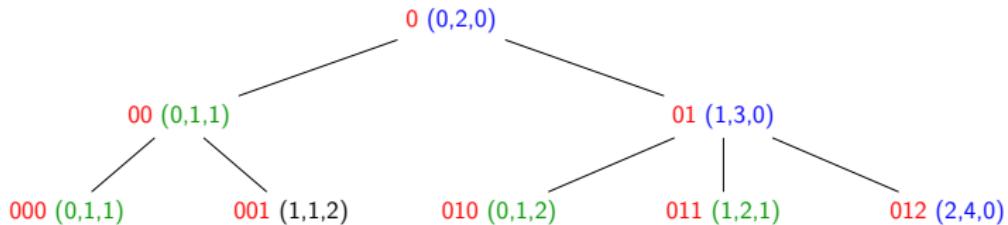


Generating Tree



- ▶ **root:** $(0, 2, 0)$
- ▶ **rules:**
 $(j - 2, j, 0) \rightarrow (j - 1, j + 1, 0), (i, i + 1, j - 1 - i)_{i=0}^{j-2}$

Generating Tree



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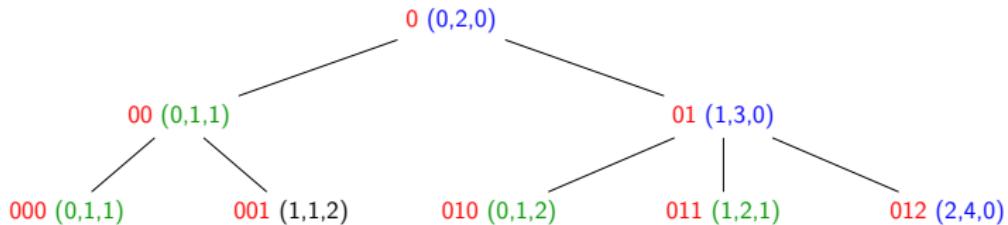
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$$(j-2, j, 0) \rightarrow (j-1, j+1, 0), (i, i+1, j-1-i)_{i=0}^{j-2}$$

$$(j-1, j, k) \rightarrow$$

$$(j-1, j, k), (i, i+1, j+k-1-i)_{i=1}^{j-2}, (j, j, i)_{i=2}^{k+1}$$

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► **root:** $(0, 2, 0)$

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$$(j, j, k) \rightarrow (j, j, k), (i, i+1, j+k-1-i)_{i=0}^{j-1}, (j, j, i)_{i=2}^k$$

Counting Nodes

- ▶ **root:** $(0, 2, 0)$

- ▶ **rules:**

$$(j-2, j, 0) \rightarrow (j-1, j+1, 0), (i, i+1, j-1-i)_{i=0}^{j-2}$$

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$$(j, j, k) \rightarrow (j, j, k), (i, i+1, j+k-1-i)_{i=0}^{j-1}, (j, j, i)_{i=2}^k$$

Note: One **blue** node per level of tree.

Need to look at **green** and **black** nodes more closely.

Tree rules to functional equations



Valparaiso
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Summary

Define:

- $g0_{n,j,k}$ is number of (j, j, k) nodes at level n .
- $B(x, y, z) = \sum_{n \geq 3} \sum_{j \geq 1} \sum_{k \geq 2} g0_{n,j,k} x^j y^k z^n$.
- $g1_{n,j,k}$ is number of $(j - 1, j, k)$ nodes at level n .
- $G(x, y, z) = \sum_{n \geq 2} \sum_{j \geq 1} \sum_{k \geq 1} g1_{n,j,k} x^j y^k z^n$.

From the rules, we obtain:

$$\begin{aligned} B(x, y, z) = & \frac{z(1 - 2y)}{1 - y} B(x, y, z) + \frac{zy^2}{1 - y} B(x, 1, z) + \frac{xy^2 z^2}{(1 - z)(1 - yz)(1 - xz)} \\ & - \frac{zy^2}{1 - y} \left(G(x, y, z) - \frac{xyz^2}{(1 - xz)(1 - yz)} \right) \\ & + \frac{zy^2}{1 - y} \left(G(x, 1, z) - \frac{xz^2}{(1 - xz)(1 - z)} \right) \end{aligned}$$

$$\begin{aligned} G(x, y, z) = & \frac{xyz^2}{(1 - xz)(1 - yz)} + \frac{zx}{x - y} G(x, y, z) \\ & - \frac{zx}{x - y} G(y, y, z) + \frac{zx}{x - y} B(x, y, z) - \frac{zx}{x - y} B(y, y, z) \end{aligned}$$

Black node data

$A0_n$ is an $(n - 2) \times (n - 2)$ array with $g0_{n,j,k}$ in row j , column $k - 1$.

$$A0_3 = \begin{bmatrix} 1 \end{bmatrix} \quad A0_4 = \begin{bmatrix} 4 & 1 \\ 1 & 0 \end{bmatrix} \quad A0_5 = \begin{bmatrix} 14 & 6 & 1 \\ 4 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix}$$

$$A0_6 = \begin{bmatrix} 50 & 27 & 8 & 1 \\ 14 & 6 & 1 & 0 \\ 4 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix} \quad A0_7 = \begin{bmatrix} 187 & 113 & 44 & 10 & 1 \\ 50 & 27 & 8 & 1 & 0 \\ 14 & 6 & 1 & 0 & 0 \\ 4 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$A0_8 = \begin{bmatrix} 730 & 468 & 212 & 65 & 12 & 1 \\ 187 & 113 & 44 & 10 & 1 & 0 \\ 50 & 27 & 8 & 1 & 0 & 0 \\ 14 & 6 & 1 & 0 & 0 & 0 \\ 4 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

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- ▶ Let $f(z) = \frac{1-z-\sqrt{1-6z+5z^2}}{2z}$.
- ▶ Let $g(z) = \frac{16z^2(z-1)}{(1-z+\sqrt{1-6z+5z^2})^3(-1+3z+\sqrt{1-6z+5z^2})}$.
- ▶ Experimentally predict:
Column i has generating function $\frac{f(z)-1}{1-z} g(z)^{i-1}$.
 $B(x, y, z) = \frac{2xy^2z^3}{(1-xz)((1-(y+1)z)\sqrt{5z^2-6z+1}+(1-(y+3)z)(1-z))}$

Green node data

$A1_n$ is an $(n - 1) \times (n - 1)$ array with $g1_{n,j,k}$ in row j column k .

$$A1_2 = \begin{bmatrix} 1 \end{bmatrix} \quad A1_3 = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \quad A1_4 = \begin{bmatrix} 1 & 3 & 1 \\ 1 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix}$$

$$A1_5 = \begin{bmatrix} 1 & 8 & 5 & 1 \\ 1 & 3 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix} \quad A1_6 = \begin{bmatrix} 1 & 23 & 19 & 7 & 1 \\ 1 & 8 & 5 & 1 & 0 \\ 1 & 3 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$A1_7 = \begin{bmatrix} 1 & 74 & 69 & 34 & 9 & 1 \\ 1 & 23 & 19 & 7 & 1 & 0 \\ 1 & 8 & 5 & 1 & 0 & 0 \\ 1 & 3 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 \end{bmatrix} \quad A1_8 = \begin{bmatrix} 1 & 262 & 256 & 147 & 53 & 11 & 1 \\ 1 & 74 & 69 & 34 & 9 & 1 & 0 \\ 1 & 23 & 19 & 7 & 1 & 0 & 0 \\ 1 & 8 & 5 & 1 & 0 & 0 & 0 \\ 1 & 3 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

Green node data

$A1_n$ is an $(n - 1) \times (n - 1)$ array with $g1_{n,j,k}$ in row j column k .

$$A1_8 = \begin{bmatrix} 1 & 262 & 256 & 147 & 53 & 11 & 1 \\ 1 & 74 & 69 & 34 & 9 & 1 & 0 \\ 1 & 23 & 19 & 7 & 1 & 0 & 0 \\ 1 & 8 & 5 & 1 & 0 & 0 & 0 \\ 1 & 3 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

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A091698 Matrix inverse of triangle A063967. 4

1, -1, 1, 1, -3, 1, -1, 8, -5, 1, 1, -23, 19, -7, 1, -1, 74, -69, 34, -9, 1, 1, -262, 256, -147, 53, -11, 1, -1, 993, -986, 615, -265, 76, -13, 1, 1, -3943, 3935, -2571, 1235, -431, 103, -15, 1, -1, 16178, -16169, 10862, -5591, 2216, -653, 134, -17, 1, 1 ([list](#) [table](#) [graph](#) [refs](#) [listen](#) [history](#) [text](#) [internal format](#))

OFFSET 0,5

COMMENTS Riordan array $(1/(1+x), (\sqrt{1+6x+5x^2}-x-1)/(2(1+x)))$. The absolute value array is $(1/(1-x), xc(x)/(1-xc(x)))$ where $c(x)$ is the g.f. of [A000108](#). It factorizes as $(1/(1-x), x/(1-x))(1, xc(x))$. - [Paul Barry](#), Jun 10 2005

LINKS [Table of n, a\(n\) for n=0..55.](#)

EXAMPLE 1; -1; 1; -3; 1; -1; 8; -5; 1; 1; -23; 19; -7; 1; ...
Contribution from [Paul Barry](#), Apr 15 2010: (Start)

Triangle begins
1,
-1, 1,
1, -3, 1,
-1, 8, -5, 1,
1, -23, 19, -7, 1,
-1, 74, -69, 34, -9, 1,
1, -262, 256, -147, 53, -11, 1,
-1, 993, -986, 615, -265, 76, -13, 1,
1, -3943, 3935, -2571, 1235, -431, 103, -15, 1

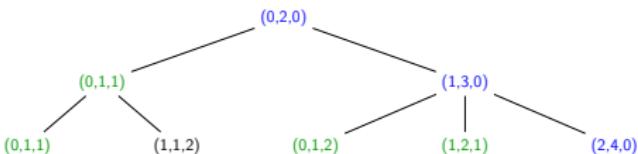
$$G(x, y, z) = \frac{2xyz^2}{(1-xz)(y\sqrt{5z^2-6z+1}+yz-2z-y+2)}$$

Recap

Goal

$$a_{0021}(n) = \sum_{\ell=0}^{n-1} \binom{n-1}{\ell} C_\ell$$

Technique: Count blue, black, and green nodes in the following generating tree:



governed by:

► **root:** $(0, 2, 0)$

► **rules:**

$$(j-2, j, 0) \rightarrow (j-1, j+1, 0), (i, i+1, j-1-i)_{i=0}^{j-2}$$

$$(j-1, j, k) \rightarrow (j-1, j, k), (i, i+1, j+k-1-i)_{i=1}^{j-2}, (j, j, i)_{i=2}^{k+1}$$

$$(j, j, k) \rightarrow (j, j, k), (i, i+1, j+k-1-i)_{i=0}^{j-1}, (j, j, i)_{i=2}^k$$

Final Enumeration

- ▶ blue nodes on level n

Generating function: $\frac{z}{1-z}$

- ▶ black nodes of type (j, j, k) on level n

Generating function:

$$B(x, y, z) = \frac{2xy^2z^3}{(1-xz)((1-(y+1)z)\sqrt{5z^2-6z+1}+(1-(y+3)z)(1-z))}$$

- ▶ green nodes of type $(j-1, j, k)$ on level n

Generating function:

$$G(x, y, z) = \frac{2xyz^2}{(1-xz)(y\sqrt{5z^2-6z+1}+yz-2z-y+2)}$$

- ▶ total nodes at level n :

$$\frac{z}{1-z} + B(1, 1, z) + G(1, 1, z) = \frac{1-z-\sqrt{5z^2-6z+1}}{2(1-z)}.$$

Summary

- ▶ There are a plethora of nice enumeration questions for pattern-avoiding ascent sequences.
- ▶ Computing $a_{0021}(n)$ completes Wilf-equivalence for patterns of length 4.
- ▶ Open: find a statistic $\text{st} : \mathcal{A}_{0021}(n) \rightarrow \ell$ so that

$$|\{a \in \mathcal{A}_{0021}(n) \mid \text{st}(a) = \ell\}| = \binom{n-1}{\ell} C_\ell.$$

References

- ▶ A. Baxter and L. Pudwell, Ascent sequences avoiding pairs of patterns, arXiv:1406.4100, submitted.
- ▶ D. Callan, T. Mansour, and M. Shattuck, Restricted ascent sequences and Catalan numbers, *Appl. Anal. Discrete Math.* **8** (2014), 288–303.
- ▶ P. Duncan and E. Steingrímsson, Pattern avoidance in ascent sequences, *Electron. J. Combin.* **18**(1) (2011), #P226 (17pp).
- ▶ T. Mansour and M. Shattuck, Some enumerative results related to ascent sequences, *Discrete Math.* **315-316** (2014), 29–41.
- ▶ L. Pudwell, Ascent sequences and the binomial convolution of Catalan numbers, arXiv:1408.6823, to appear in *Australas. J. Combin.*.

References

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Thanks for listening!