A58929 Scan Chae, Palmer, RWR all 65 segs

CHAE

A58929

Computing the number of Claw-free Cubic Graphs

with given Connectivity

Gab-Byung Chae chaegabb@pilot.msu.edu

Edgar M. Palmer palmer@math.msu.edu Department of Mathematics Michigan State University East Lansing, MI, 48823

Robert W. Robinson
rwr@pollux.cs.uga.edu
Computer science Dept.
415 GSRC
University of Georgia
Athens, GA 30602-7404
October 31, 2000

#### Abstract

We use exponential generating functions to count claw-free cubic graphs with given connectivity. Tables are provided for connectivity 1, 2 and 3.

#### 1 Introduction

A claw-free cubic graph G is a cubic graph which contains no induced subgraph isomorphic to  $K_{1,3}$ . Therefore these are precisely the cubic graphs whose vertices all

belong to triangles. For convenience we refer to them as *cfc's*. For a study of the history of the enumeration of regular graphs the reader can consult Gropp [Gr92] and for more recent work on enumerating cubic and claw-free cubics see [PaRR0x], and [CPR0x].

 $\langle$ 

In [PaRR0x] Palmer, Read, and Robinson studied claw-free cubic graphs and computed the number of claw-free cubic graphs with up to 52 vertices. Moreover, their paper contained a partial differential equation for the exponential generating function(egf) of labeled, general cubic graphs. In [CPR0x] this equation was used to derive recurrence relations for general cubic graphs with a specified number of multiple edges and loops by connectedness. There is another relevant paper with enumeration formulae for cfc's, namely [MPaRR0x]. Combining results of these papers makes it possible to count claw-free cubic graphs with given connectivity.

In the present paper, we will follow the terminology and method in [PaRR0x] to find the number of k-connected claw-free cubic graphs where k = 1, 2, and 3. In a claw-free cubic graph, every vertex belongs to a triangle. So the maximum number of triangles in which a vertex may lie is 3. Clearly, a vertex lies in 3 triangles if and only if it is a vertex of  $K_4$ . A diamond in a cfc is an induced subgraph isomorphic to  $K_4 - e$ . A vertex lies in exactly two triangles if and only if it is one of the vertices of degree 3 in a diamond. A string of diamonds is a induced subgraph in which diamonds are adjacent in series. A ring of diamonds is a connected component in which every vertex belongs to a diamond.

In [PaRR0x] there are two important operations which convert general cubic graphs to claw-free cubic graphs. One of them is the *expansion* operation which converts an edge of a general cubic graph to a string of diamonds. The inverse operation to *expansion* is called *reduction*. The other operation is *dilation* which inflates a vertex of a general cubic graph to a triangle. The inverse operation to *dilation* is called *contraction*. Consider a cfc with no component isomorphic to  $K_4$  or a ring of diamonds. The reduction operation applied to all strings of diamonds in this cfc

results in a general cubic graph with no loops but possibly some double edges. The two vertices of such a double edge are mutually adjacent to a third vertex. These vertices constitute a trumpet. Still, every vertex must belong to a triangle. Note that if e is an edge caused by a reduction, then such a triangle must be part of a trumpet in which e is a multiple edge. Now the contraction operation completes the conversion from claw-free cubic graphs to general cubic graphs.

For general graph theoretic terminology and notation we follow [CL96] and the basic knowledge of labeled enumeration techniques can be found in [HPa73].

# 2 Connected claw-free cubic graphs

We define

$$H(z) = \sum_{n=0}^{\infty} h_n \frac{z^n}{(2n)!}.$$

where  $h_n$  is the number of labeled claw-free cubic graphs on 2n vertices. Then  $H(z^2)$  is the exponential generating function for these graphs. By applying expansion and dilation operations, Palmer, Read, and Robinson [PaRR0x] derived the following a differential equation whose formal power series solution is the egf for labeled claw-free cubic graphs on 2n vertices:

$$0 = (144z^{8} + 288z^{7} - 576z^{4})H''(z)$$

$$+ (-36z^{10} - 96z^{9} + 24z^{8} + 144z^{7} + 576z^{6} + 384z^{5}$$

$$- 576z^{4} - 2880z^{3} - 576z^{2} + 1152)H'(z)$$

$$+ (-15z^{11} - 74z^{10} - 130z^{9} - 96z^{8} + 144z^{7} + 368z^{6} + 336z^{5} - 288z^{4}$$

$$- 240z^{3} - 288z^{2} - 96z)H(z).$$

$$(1)$$

Equation (1) can be converted to a differential equation whose formal solution is the egf for the number of connected, claw-free cubic graphs by the substitution

$$H(z) = e^{H_1(z)},\tag{2}$$

enor!

 $h_1(6) = 59875200$ 

Table 1: Boundary conditions.

 $h_1(1) = 0$   $h_1(7) = 13621608000$ 

 $h_1(2) = \widehat{12}$   $h_1(8) = 8009505504000$  $h_1(3) = 60$   $h_1(9) = 3123380227968000$ 

 $h_1(4) = 2520$   $h_1(10) = 1832279324908032000$ 

 $h_1(5) = 453600$   $h_1(11) = 2054813830468439040000$ 

 $h_1(12) = 1665031453088810526720000$ 

where  $H_1(z^2)$  is the egf for connected, claw-free cubic graphs, i.e.

$$H_1(z) = \sum_{n=0}^{\infty} h_1(n) \frac{z^n}{(2n)!},$$

and  $h_1(n)$  is the number of connected, labeled cfc graphs with 2n vertices. After substitution in equation (1) of H(z) and its derivatives from equation (2), we have the following differential equation for  $H_1(z)$ :

$$0 = (144z^{8} + 288z^{7} - 576z^{4})H'_{1}(z)H'_{1}(z)$$

$$+ (144z^{8} + 288z^{7} - 576z^{4})H''_{1}(z)$$

$$+ (-36z^{10} - 96z^{9} + 24z^{8} + 144z^{7} + 576z^{6} + 384z^{5}$$

$$- 576z^{4} - 2880z^{3} - 576z^{2} + 1152)H'_{1}(z)$$

$$+ (-15z^{11} - 74z^{10} - 130z^{9} - 96z^{8} + 144z^{7} + 368z^{6} + 336z^{5} - 288z^{4}$$

$$- 240z^{3} - 288z^{2} - 96z).$$
(3)

The recurrence relation for the number of connected, claw-free cubic graphs can be found by extracting the coefficient of  $\frac{z^n}{(2n)!}$  from both sides of (3). The relation is supported by the boundary conditions in Table 1:

For  $n \ge 13$ , we have:

$$h_{1}(n) = -144 \frac{(2n)!}{1152n} \sum_{k=1}^{n-7} \frac{kh_{1}(k)(n-7-k)h_{1}(n-7-k)}{(2k)!(2n-14-2k)!}$$

$$-288 \frac{(2n)!}{1152n} \sum_{k=1}^{n-6} \frac{kh_{1}(k)(n-6-k)h_{1}(n-6-k)}{(2k)!(2n-12-2k)!}$$

$$+576 \frac{(2n)!}{1152n} \sum_{k=1}^{n-3} \frac{kh_{1}(k)(n-3-k)h_{1}(n-3-k)}{(2k)!(2n-6-2k)!}$$

$$-144 \frac{(2n)!}{(1152n)(2n-14)!}(n-7)(n-8)h_{1}(n-7)$$

$$-288 \frac{(2n)!}{(1152n)(2n-12)!}(n-6)(n-7)h_{1}(n-6)$$

$$+576 \frac{(2n)!}{(1152n)(2n-6)!}(n-3)(n-4)h_{1}(n-3)$$

$$+36 \frac{(2n)!}{(1152n)(2n-20)!}(n-10)h_{1}(n-10)$$

$$+96 \frac{(2n)!}{(1152n)(2n-18)!}(n-9)h_{1}(n-9)$$

$$-24 \frac{(2n)!}{(1152n)(2n-16)!}(n-8)h_{1}(n-8)$$

$$-144 \frac{(2n)!}{(1152n)(2n-14)!}(n-7)h_{1}(n-7)$$

$$-576 \frac{(2n)!}{(1152n)(2n-12)!}(n-6)h_{1}(n-6)$$

$$-384 \frac{(2n)!}{(1152n)(2n-10)!}(n-5)h_{1}(n-5)$$

$$+576 \frac{(2n)!}{(1152n)(2n-8)!}(n-4)h_{1}(n-3)$$

$$+2880 \frac{(2n)!}{(1152n)(2n-6)!}(n-3)h_{1}(n-3)$$

$$+576 \frac{(2n)!}{(1152n)(2n-6)!}(n-2)h_{1}(n-2).$$

By applying *Mathematica* to this recurrence relation, we calculated the numbers of 1-connected cfc graphs shown in Table 2. Actually the boundary conditions are not found only by the equation (3) directly. It just gives us partial values of them. In order to find exact values, we have to add the contribution from the constants in

equation (3) to the values which come from the output of above recurrence relation (4) for n up to 12.

For example, when n = 12 the number 15\*24!/(1152\*12) = 673229602575129600000 comes from the constants in equation (3) and the value 1664358223486235397120000 comes from the recurrence relation (4) by using the previous boundary values for n < 12. The sum of these two numbers gives us the boundary condition  $h_1(12) = 1665031453088810526720000$ .

### 3 2-connected Claw-free cubic graphs

A 2-connected general cubic graph can be converted to a 2-connected cfc graph by the expansion operation, which converts an edge of a general cubic graph to a string of diamonds, and the dilation operation, which inflates a vertex of a general cubic graph to a triangle. The smallest 2-connected general cubic graph is  $K_4$ . By the dilation operation  $K_4$  is converted to a 2-connected cfc graph of order 12. Therefore these operations produce 2-connected cfc graphs of order at least 12. However, there are 2-connected cfc graphs that can not be produced in this way. In fact there are three types of such graphs:

- (a) The triangular prism of order 2n = 6.
- (b) The rings of two or more diamonds of order  $2n \geq 8$ .
- (c) The graphs of order  $2n \ge 10$  obtained by expanding with diamonds the edges of the triangular prism that do not belong to a triangle.

Next we need the egf's for these three families. Of course,  $K_4$  is counted by  $z^2/4!$ .

The rings of diamonds are counted by
$$= \frac{1}{163} + \frac{1}{48} = \frac{3}{128}$$

$$= \frac{1}{128} = \frac{3}{128} + \frac{1}{128} = \frac{3}{128} =$$

$$\sum_{m=2}^{\infty} \frac{z^{2m}}{2m2^m} = -\frac{z^2}{4} - \frac{1}{2}ln(1 - z^2/2). \tag{5}$$

And 2-connected cfc graphs obtained by expanding with diamonds the edges of

Table 2: Number of 1-connected cfc graphs with  $2n \le 60$ .

$\overline{n}$	$h_1(n)$	
1	0	
2	1	
3	60	
4		
5	2520 453600 59875200 13621608000	
6	59875200	
7	13621608000	
8	8009505504000	
9	3123380227968000	
10	1832279324908032000	
11	2054813830468439040000	
12	1665031453088810526720000	
13	1925086583971531588608000000	
14	3552833935369312965955584000000	
15	5046746501122027301952608256000000	
16	9861817424365745824355502612480000000	
17	27365975784025025428617030645350400000000	
18	61323963707903030791402423349300428800000000	
19	183552463622453002911375211047071799705600000000	
20	720052647634369560568722076458423100794470400000000	
21	2368360586025317757755816851785727694336557056000000000	
22	101467845834916695851867212426301854404400565452800000000000	
23	53795556323350118084055188978516784012472039393198080000000000000	
24	2464244707795626835290012843752032357333546137959163494400000000000	
25	$\left \begin{array}{c}1438454000284443072212393572236725837273648853200978470502400000000000\end{array}\right $	
26	$\left  995856269134237859148875229290819201228673905380011308954419200000000000000000000000000000000000$	
	00	
27	$\left  610405473680692784571394583659731682395837369994603392192447774720000000 \right $	
	000000	
28	46763587708656896789988944665808439580376015582411036159102127439872000	
	000000000	
29	41141634055363776663530319633729522624244890505593059232232304816003481	
0.0	600000000000	
30	32697533978263237827530129558757067499130106959793034875270154251141203	
	8860800000000000	

the triangular prism that do not belong to a triangle has egf

$$z^3b^3/12$$

where

$$b(z) = \sum_{k=0}^{\infty} \left(\frac{z^2}{2}\right)^k = (1 - z^2/2)^{-1}$$
 (6)

is the egf of the strings of diamonds.

Let  $\Phi(z^2)$  be the egf of all three types above. Then

$$\Phi(z) = z^2/24 + \sum_{m=2}^{\infty} \frac{z^{2m}}{2m2^m} + \sum_{k=0}^{\infty} \frac{(k+1)(k+2)}{24} \frac{z^{2k+3}}{2^k}.$$

Let  $G_2(x,y)$  be the egf

$$G_2(x,y) = \sum_{s,d} g_2(2m,d) \frac{x^s y^d}{(2n)!}$$

where  $g_2(2m, d)$  is the number of 2-connected labeled general cubic graphs of order 2m with s single edges, d double edges and  $2m = \frac{2s+4d}{3}$ . We define  $f_2(2n, d)$  to be the number of cfc's of order 2n built from 2-connected general cubic graphs with s single edges, d double edges and no loops by dilating vertices and expanding edges. Then we have the following formula which is simpler then the one in [MPaRR0x] because we do not have loops in 2-connected cfc's.

**Lemma 3.1** For fixed n, d, we have

$$f_2(2n,d) = \sum_{m,j} g_2(2m,d) \binom{2n}{6m} \frac{\binom{6m}{3,\cdots,3}}{(2m)!} (3!)^{2m-2d} (3^2 \cdot 2)^d \left(\binom{3m}{j}\right) \binom{4j}{4,\cdots,4} (12)^j,$$
(7)

where j is the number of diamonds and  $2n \geq 12$ .

**Proof.** Suppose G is a 2-connected labeled general cubic graphs counted by  $g_2(2m,d)$ . First we choose 6m labels from 2n available and arrange them in 2m unordered groups of three each for triangles. Then

$$\binom{2n}{6m} \frac{\binom{6m}{3,\cdots,3}}{(2m)!}$$

is the number of ways to do this. In triangles, the number of ways to label the vertices according to the adjacencies is

$$(3!)^{2m-2d}(3^2 \cdot 2)^d.$$

Since there are 3m original edges in the 2-connected labeled general cubic graph G, they can be expanded by j diamonds using combinations with repetition we find that the number of ways to do this is

$$\left( \begin{pmatrix} 3m \\ j \end{pmatrix} \right) = \begin{pmatrix} 3m+j-1 \\ j \end{pmatrix}.$$

The number of ways to arrange the remaining labels for the diamonds and the the number of ways to assign labels to individual diamonds is

$$\binom{4j}{4,\cdots,4}(4!/2)^j.$$

Note that the number of 2-connected cfc graphs which can be obtained by (7) depends on the number of double edges in 2-connected labeled general cubic graphs  $g_2(2m, d)$  which were already computed in [CPR0x]. Define

$$B_2(z^2) = \sum_{n=0}^{\infty} b_2(2n) \frac{z^{2n}}{(2n)!}$$

be the egf of 2-connected cfc graphs which can be obtained by (7), then  $b_2(2n) = \sum_d f_2(2n, d)$ . Let

$$H_2(z^2) = \sum_{n=0}^{\infty} h_2(n) \frac{z^{2n}}{(2n)!}$$

be the egf of 2-connected cfc graphs. Then we have

$$H_2(z^2) = B_2(z^2) + \Phi(z^2).$$

But the computing the number of 2-connected cfc graphs by using this egf is not quite simple. To get the  $b_2(2n)$ , we need to compute the number  $f_2(2n,d)$  and sum them up according to the number of vertices 2n. And then extract the coefficients of three egf's in  $\Phi(z^2)$ . By adding, finally, the above numbers from each egf's according to the number of vertices, we can have the numbers of 2-connected cfc graphs as in Table 3.

Table 3: Number of 2-connected cfc graphs with  $2n \leq 36$ .

n	$h_1(n)$	
2	1	
3	60	
4	60 2520 A 5 8929	
5	453600	
6	59875200	
7	10897286400	
8	6701831136000	
9	2623194782208000	
10	1338096104497152000	
11	1633313557551836160000	
12	1324107982344764897280000	
13	1408369399403068118016000000	
14	2818005386051236981856256000000	
15	3984871608553561924638375936000000	
16	7418092561827244386962686894080000000	
17	22027134615845465196052794703872000000000	
18	49003622223231250364949254126429798400000000	

### 4 3-connected Claw-free cubic graphs

Let  $G_3(x,y)$  be the egf

$$G_3(x) = \sum_s g_3(2m) \frac{x^s}{(2m)!}$$

where  $g_3(2m)$  is the number of 3-connected labeled general cubic graphs of order 2m with s single edges and  $2m = \frac{2s}{3}$ . Then we define  $f_3(2n)$  to be the number of cfc's of order 2n built by dilating vertices in general cubic graphs with s single edges, no double edges and no loops. Suppose G is a 3-connected labeled general cubic graph counted by  $g_3(2m)$ . Then the contribution of G to  $f_3(2n)$ , with  $(2m) \cdot 3 = 2n$ , is determined by by arranging the 2n labels in 2m unordered groups of three vertices each for triangles. Here is the simple relationship between  $f_3(2n)$  and  $g_3(2m)$ .

Table 4: Number of 3-connected cfc graphs with  $2n \leq 90$ .

n	$h_1(n)$
2	1
3	60 19958400 A 58 9 30
6	19958400
9	622452999168000
12	258520167388849766400000
15	675289572271869736778268672000000
18	7393367369949286697176489031997849600000000
21	2627800504609683185243971405741688045646643200000000000
24	25427675465852111040703353545981158863084030467978035200000000000
27	588899571830694942000264105108811607070150958832700604777758720000000
	0000000
30	295804325421925626330882127682606558611431840588595256206808510098367
	840256000000000000
33	298089015291900801910918687858981579022518884435603158506899447288503
	7078284530089984000000000000000
36	565677772026602700573118887454300022325482015168187354914748390689609
	116481451145434282929946624000000000000000000000000000000
39	191803897837508699578197474721718346802046111134993250012544269312632
	860958978793819181561437812810831626240000000000000000000000000000000000
42	111159153925422985672391611050830189648350405191904965108689674793493
	964995720171415053312805183121693748717674299392000000000000000000000000000000000
45	106005759806440161267490042030591700071014694671953768589449762663365
	484968378587742696877946003511137006557545399304494112361676800000000
	00000000000

#### **Lemma 4.1** For fixed n, we have

$$f_3(2n) = g_3(2m)\frac{(6m)!}{(2m)!},\tag{8}$$

where 2n = 6m.

The smallest 3-connected general cubic graph is  $K_4$ . Therefore the dilation operation produces 3-connected cfc graphs of order at least 12. But it will produce every 3-connected cfc graph except the triangular prism of order 2n = 6. The numbers  $g_3(2m)$ 

Table 5: Number of unlabeled and labeled cfc with  $\kappa(G) = 1$ 

2n	# of unlabeled cfc with $\kappa(G) = 1$	# of labeled cfc with $\kappa(G) = 1$
14	1	2724321600
16	1	1307674368000
18	3 A58932	500185445760000
20	5	494183220410880000
22	11	421500272916602880000
24	20	340923470744045629440000

can be found by using Wormald's recurrence relation [W79c]. We have

$$g_3(2m) = (2m)! \frac{r(m)}{3m2^m}$$

where

$$r(m) = (3m-2)(r(m-1) + \sum_{i=2}^{m-2} r(i)r(m-i).$$

We used this method to compute the numbers  $g_3(2m)$ , *i.e* the number of 3-connected cfc graphs, shown in 4. This complete the enumeration of cfc graphs with given connectivity.

# 5 Conclusion

The values of  $h_1(n)$  and  $h_3(n)$  were checked for  $n \leq 12$  by calculating the order of the automorphism graphs of the small connected cfc's.

The numbers  $h_2(n)$  were also checked for small value of n by finding the diagrams of the unlabeled cfc's with connectivity 1. These are the graphs that contribute to the difference between  $h_1(n)$  and  $h_2(n)$ . For example, there are 20 unlabeled cfc's with 2n = 24 vertices and  $\kappa(G) = 1$ . And the number of ways to label these is 340923470744045629440000. And then we compare this to the number which is the

difference between the  $h_1(12)$  and  $h_2(12)$  which are found in the Tables 2 and 3, respectively as follows:

$$340923470744045629440000 = h_1(12) - h_2(12)$$

The numbers on the right side of the Table 5 are exactly the differences between  $h_1(n)$  and  $h_2(n)$ .

Finally, we note that almost 80% of cfc's are 2-connected when the number of vertices is 34 or 36. This is consistent with the observation in [MPaRR0x] that almost all cfc's are 2-connected.

#### References

- [CPR0x] G. Chae, E.M. Palmer, & R.W. Robinson, Computing the number of labeled general cubic graphs, to appear.
- [CL96] G. Chartrand and L. Lesniak, *Graphs and Digraphs*, Chapman & Hall, 2-6 Boundary Row, London (1996).
- [Gr92] Harald Gropp, Enumeration of regular graphs 100 years ago, Discrete Mathematics 101 (1992) 73-85.
- [HPa73] F. Harary and E.M. Palmer, *Graphical Enumeration*, Academic, New York (1973).
- [MPaRR0x] B.D. McKay, E.M. Palmer, R.C. Read, & R.W. Robinson, The asymptotic number of claw-free cubic graphs, *Proc. Eighth QuadrennialInternat. Conf. on Graph Theory, Combinatorics, Algorithms, and Applications* (Kalamazoo, 1996), to appear.
- [PaRR0x] E.M. Palmer, R.C. Read, & R.W. Robinson, Counting claw-free cubic graphs, to appear.

- [RW80] R.C. Read and N.C. Wormald, Number of labeled 4-regular graphs, *Journal of Graph Theory*, 4 (1980) 203-212.
- [W79c] N.C. Wormald, Enumeration of labelled graphs II: Cubic graphs with a given connectivity, J. London Math.Soc. (2) 20 (1979) 1-7.