# **15-396** Science of teh Interwebs

# **Preliminaries of Game Theory**

Lecture 8 (October 4, 2011)





Payoffs

Strategies for Player 1:  $S_1 = \{WiFi, Wired\}$ Strategies for Player 2:  $S_2 = \{WiFi, Wired\}$  $S = \{(WiFi, Wired), (WiFi, WiFi), (Wired, WiFi), (Wired, Wired)\}$  We assume that everything a player cares about is summarized in the player's payoff

We also assume that each player knows everything about the game

### **Prisoner's Dilemma**



### **Best Responses**

A strategy  $s_1^*$  is a best response by player 1 to a strategy  $s_2$  for player 2 if

 $\pi_1(s_1^*, s_2) \ge \pi_1(s_1, s_2)$ 

for all strategies  $s_1 \in S_1$ .



If Suspect 2 does not confess, then confessing is a best response for Suspect 1

# **Dominant Strategy**

A strategy  $s_1^*$  is a Dominant Strategy for player 1 if  $s_1^*$  is a Best Response to every possible strategy for player 2.



# Confessing is a dominant strategy for both Suspects!



#### I is a dominant strategy for both players



#### (L,H) will be played



# Neither player has a dominant strategy

# Nash Equilibrium

A pair of strategies  $(s_1^*, s_2^*)$  is in Nash Equilibrium if  $s_1^*$  is a Best Response by player 1 to  $s_2^*$ , and  $s_2^*$ is a Best Response by player 2 to  $s_1^*$ .



# **Coordination Game**

#### Player 2



Nash Equilibria: (L,L), (R,R)



#### Nash Equilibria: (D,H), (H,D)

# **Matching Pennies**

Player 2



#### No pure Nash Equilibria Exist!

## **Randomized Strategies**



Player 1 picks H with probability p and Player 2 picks H with probability q



E[Payoff for P1 doing H] = (-1)q + (+1)(1-q) = 1-2qE[Payoff for P1 doing T] = (+1)q + (-1)(1-q) = 2q-1Player 1 will choose H if 1-2q > 2q-1. i.e., if q < 1/2Player 1 will choose T if 1-2q < 2q-1. i.e., if q > 1/2

#### We say that (p\*,q\*) is a mixed strategy Nash Equilibrium if p\* is a best response by player 1 to q\* and q\* is a best response by player 2 to p\*



(p=2/3,q=3/4) is an equilibrium!

Player 1 is only willing to randomize if the expected payoffs of U and D are equal: q+4(1-q)=2q+(1-q), so q=3/4





But (U,L) gives each player a payoff of 1, whereas (D,R) gives them 2.

Nash Equilibrium not always socially optimal

